Database Indexes, Triggers and Transactions

Databases and Web Applications Laboratory (LBAW) Bachelor in Informatics Engineering and Computation (L.EIC)

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Outline

- → Database Specification (EBD) development (A6)
 - → Indexes
 - → Triggers
 - → Transactions
 - → Database Population

PostgreSQL setup

LBAW Plan

→ Plan: <u>https://web.fe.up.pt/~ssn/wiki/teach/lbaw</u>

- \rightarrow 5th week of classes;
- \rightarrow Continue development on the second component (EBD);
- \rightarrow Lab classes:
 - \rightarrow continue work on component (EBD);
 - \rightarrow work on the relational schema (A5).
- → Monitor sessions: Wednesday, at 15h, online
 - → PostgreSQL setup and use.



Database Specification (EBD) Development

Database Specification (EBD) Component

- The EBD component groups the artefacts to be made by the development team in order to support the storage and retrieval requirements identified in the requirements specification.
- \rightarrow It consists of three artefacts:
 - → A4: Conceptual Data Model
 - → A5: Relational Schema, Validation and Schema Refinement
 - → A6: Indexes, Triggers, Transactions and Database Population

→ https://web.fe.up.pt/~ssn/wiki/teach/lbaw/202122/artefacts





A5. Relational Schema, Validation and Schema Refinement

- The A5 artifact contains the Relational Schema obtained by mapping from the Conceptual Data Model.
- The Relational Schema includes each relation schema, attributes, domains, primary keys, foreign keys and other integrity rules: UNIQUE, DEFAULT, NOT NULL, CHECK.
- \rightarrow Relation schemas are specified in the compact notation.
- \rightarrow In addition to this representation, the relational schema is also presented in SQL as an annex.
- To validate the Relational Schema obtained from the Conceptual Model, all functional dependencies are identified and the normalization of all relation schemas is accomplished.
- → Should it be necessary, in case the scheme is not in the Boyce–Codd Normal Form (BCNF), the relational schema is refined using normalization.







A5. Relational Schema Compact Notation

- \rightarrow Relation schemas are specified in the compact notation:
- \rightarrow table1(<u>id</u>, attribute NN) table2(id, attribute \rightarrow Table1 NN) table3(i<u>d1, id2 \rightarrow Table2, attribute UK NN)</u> table4((id1, id2) \rightarrow Table3, id3, attribute)
- \rightarrow Primary keys are underlined. UK means UNIQUE and NN means NOT NULL.
- → Today DATE DEFAULT CURRENT_DATE Priority ENUM ('High', 'Medium', 'Low')
- \rightarrow In PostgreSQL use lower case and the "snake_case" convention.



 \rightarrow The specification of additional domains can also be made in a compact form, using the notation:



A5. Relational Schema Mapping

Summary of Mapping Rules from **Logical UML Models to Relational Schemas**

Translated from:

UML – Metodologias e Ferramentas CASE, Vol. 1, 2ª Edição, pp. 314-315 Alberto Silva e Carlos Videira, Centro Atlântico (2005)

Rule 1	Classes are mapped into relation schemas
Rule 2	Class attributes are mapped to attributes o
Rule 3	Operations of classes are generally not map and executed in the global context of the da
Rule 4	Objects are mapped into tuples of one or m
Rule 5	Each object is uniquely identified. If the identification of an object is defined e or more attributes, this attribute is mapped Otherwise, we assume implicitly that the o name of the relation and common suffix (e.
Rule 6:	The mapping of many-to-many association together as primary key, and individually a involved.
Rule 7:	The mapping of one-to-many associations in class that has the constraint "many", of a fo

Alberto Manuel Rodrigues da Silva e Carlos Alberto Escaleira Videira, UML, metodologias e ferramentas CASE, 2ª Edição, Volume 1, Centro Atlântico Editora, Maio 2005.



of relations.

pped. They can nevertheless be mapped to *stored procedures,* stored atabase involved.

nore relations.

explicitly by the OID (object identifier) stereotype, associated with one d to primary key in the relation schema.

corresponding primary key is derived from a new attribute with the .g. "PK", "ID").

is involves the creation of a new relation schema, with attributes acting as foreign key for each of the schemas derived from the classes

involves the introduction, in the relation schema corresponding to the oreign key attribute for the other schema.





A5. Relational Schema Mapping

Rule 8:	The mapping of one-to-one associations has attributes of the classes involved in one con corresponding schema and choose one of th attribute for the other schema. This attribut
Rule 9:	Association navigability in general has no in associations, when they are complemented should include the foreign key attribute.
Rule 10:	Aggregation and composition associations h to the definition of constraints cascade ("CA
Rule 11:	The mapping of generalization associations The first solution consists in crushing the his superclass. This solution is appropriate whe and/or when the semantics of their identified The second solution is to consider only sche the super-class in these schemas; in particu The third solution is to consider all the sche of connected schemas and maintained at the advantage of avoiding duplication of inform information by various schemas, and might data by requiring the execution of various je

Alberto Manuel Rodrigues da Silva e Carlos Alberto Escaleira Videira, UML, metodologias e ferramentas CASE, 2ª Edição, Volume 1, Centro Atlântico Editora, Maio 2005.

review

s in general two solutions. The first corresponds to the fusion of the mmon schema. The second solution is to map each of the classes in the ne schemas as the most suitable for the introduction of a foreign key te should also be defined as unique within that schema.

mpact on the mapping process. The exception lies in one-to-one with navigation cues it helps in the selection of the schema that

have a minimal impact on the mapping process, which may correspond SCADE") in changing operations and/or removal of tuples.

in general presents three solutions.

ierarchy of classes in a single schema corresponding to the original en there is a significant distinction in the structure of sub-classes cation is not strong.

emas corresponding to the sub-classes and duplicate the attributes of alar it works if the super-class is defined as abstract.

emas corresponding to all classes of the hierarchy, resulting in a mesh e expense of referential integrity rules. This solution has the ation among different schemas, but suggests a dispersion of involve a performance penalty in query operations or updating of join operations (i.e. "JOIN") and/or validation of referential integrity.





Mapping Generalizations



```
# Superclass approach
```

```
media(id, type CHK {book, cd, dvd} ...)
```

ER approach

```
media(<u>id</u>, ...)
book(id->media, ...)
cd(<u>id->media</u>, ...)
dvd(<u>id->media</u>, ...)
```

Object Oriented

book(id, [media attributes], ...) cd(id, [media attributes], ...) dvd(id, [media attributes], ...)



A6. Indexes, Triggers, Transactions and Database Population

- → This artefact contains the physical schema of the database,
 - \rightarrow the identification and characterization of the indexes,
 - → the support of data integrity rules with triggers,
 - → the definition of the database user-defined functions,
 - \rightarrow and the identification and characterization of the database transactions.
- This artefact also includes the complete database creation script, including all SQL code necessary to define all integrity constraints, indexes, triggers and transactions.
- → Also, the database creation script and the database population script should be included as separate elements.

A6. Indexes

- of tuples in each relation.
- \rightarrow **Performance indexes** are applied to improve the performance of select queries.
 - impact on the performance of the application.
- GIN), and also if clustering is recommended. As a last resource, controlled redundancy may be introduced (de-normalisation).
- namely all necessary configurations, indexes definitions and other relevant details.

> The workload is a study of the predicted system load, including an estimate on the number and growth

→ At most, three performance indexes can be proposed, identifying the ones that have the biggest

-> For each proposed index, it is necessary to indicate and justify the type chosen (B-tree, Hash, GiST,

The system being developed must provide full-text search features supported by PostgreSQL. Thus, it is necessary to specify the fields where full-text search will be available and the associated setup,







Indexes in PostgreSQL (1)

- back-of-the-book index).
- - \rightarrow Without indexes, tables are usually sequentially scanned to find the matching entry.
 - \rightarrow With indexes, the number of steps to find the matching records can be drastically reduced.
- \rightarrow Two main types:
 - range search in log time.
 - operator is used (no sorting or ranges).

→ Indexes are secondary data structures used to improve data access (useful metaphor - the alphabetical

 \rightarrow Finding and retrieving specific rows is much faster with indexes, but they add an overhead to the execution.

→ B-tree indexes: use a tree-like data structure that maintains data sorted and allow for search, order,

→ Hash indexes: use a hash-function to map keys to values; are only considered when an equality



Indexes in PostgreSQL (2)

- \rightarrow Indexes can be created for more than one attribute (multicolumn).
 - \rightarrow CREATE INDEX name ON table (a, b);
- \rightarrow Indexes can also be created for expressions.
 - \rightarrow SELECT * FROM test1 WHERE lower(col1) = 'value';
 - → CREATE INDEX test1_lower_col1_idx ON test1 (lower(col1));
- \rightarrow Index usage can be analyzed with the EXPLAIN command.

 \rightarrow Work when searching for both attributes simultaneously or just a. Not just b.



Unique Indexes

- → Indexes can also be used to enforce uniqueness of a column's value, or the uniqueness of the combined values of more than one column.
- → CREATE UNIQUE INDEX name ON table (column [, ...]);
- → When an index is declared unique, multiple table rows with equal indexed values are not allowed. Null values are not considered equal.
- PostgreSQL automatically creates a unique index when
 - → a unique constraint or
 - \rightarrow primary key is defined for a table.



Clustering

- → Clustering a table results in the physical re-ordering of data in disk based on the index information. To cluster a table, an index must already be defined.
- Clustering is a one-time operation: when the table is subsequently updated, the changes are not clustered.
- → If needed, clustering can be set to run periodically using cron.
- → Clustering will help when multiple records are read together and an index can group them. It will be irrelevant when single rows are randomly accessed in a table.



Cardinality

- \rightarrow The uniqueness of data values contained in a particular column.
- \rightarrow The lower the cardinality, the more duplicate values in the column.
- \rightarrow Examples:
 - → high cardinality primary key
 - → medium cardinality last name in a customer table
 - → low cardinality boolean column
- used.

- Cardinality is used by the PostgreSQL planner, amongst other statistics, to estimate the number of rows returned by a WHERE clause. This is then used to decide if, and what, indexes should be





Full Text Search

- → How can you search for a work in text fields? And multiple words?
- → Using the LIKE operator is not feasible
 - \rightarrow There is no linguistic support (e.g. singular / plural).
 - \rightarrow No ranking is provided, only a set of results.
 - \rightarrow Multiple words search is not supported.
 - \rightarrow There is no index support.
- \rightarrow It is necessary to index each work individually.
- \rightarrow This is called *full text search*, or simply *text search*, on PostgreSQL.

 \rightarrow Key first step — define what is a document in our search system and what information is relevant.



tsvector lype

- \rightarrow Text are broken into lexemes, a normalized representation of words, e.g. normalization includes converting to lowercase, identifying the stem, etc.
- → The tsvector data type is used to store distinct lexemes.
 - SELECT to tsvector('english', 'The quick brown fox jumps over the lazy dog')
 - → 'brown':3 'dog':9 'fox':4 'jump':5 'lazi':8 'quick':2
- \rightarrow The function to tsvectors returns a tsvector with duplicates removed, stop words removed, and the number of position of each lexeme recorded.



tsqueries Type

- \rightarrow Queries, i.e. searches, are represented as tsqueries.
- The to_tsquery and plainto_tsquery functions convert a text query to a tsquery, a structure optimized for searching tsvectors.
- → SELECT plainto_tsquery('english','sail boats');
 - → 'sail' & 'boat'
- → SELECT plainto_tsquery('portuguese', 'o velho barco');
 - → 'velh' & 'barc'



Matching tsqueries to tsvectors

- \rightarrow The @@ operator is used to assert if a tsvector matches a tsquery:
 - → SELECT to_tsvector('portuguese','o velho barco') @@ plainto_tsquery('portuguese', 'barca');

→ †

→ SELECT to_tsvector('portuguese', 'o velho barco') @@ plainto_tsquery('portuguese','carro');

→ f

→ SELECT title FROM posts WHERE to_tsvector('english', title || ' ' || body) @@ plainto_tsquery('english', 'jumping dog');

·		

FTS Weights

- Sometimes we want to give more importance to some specific fields.
- \rightarrow We can use the setweight to attach a weight to a certain tsvector.
- \rightarrow Weights go from 'A' (more important) to 'D' (less important).
- → SELECT

setweight(to_tsvector('english', 'The quick brown fox jumps over the lazy dog'), 'A') setweight(to_tsvector('english', 'An English language pangram. A sentence that contains all of the letters of the alphabet.'), 'B')

→ 'alphabet':24B 'brown':3A 'contain':17B 'dog':9A 'english':11B 'fox':4A 'jump':5A 'languag':12B 'lazi':8A 'letter':21B 'pangram':13B' quick':2A 'sentenc':15B



Ranking FTS Results

- PostgreSQL provides two predefined ranking functions, which take into account lexical, proximity, and structural information:
 - → how often the query terms appear in the document;
 - → how close together the terms are in the document;
 - → how important is the part of the document where they occur.
- Different applications might require additional information for ranking, e.g., document modification time. The built-in ranking functions are only examples.

Ranking FTS Results

- between a tsquery and tsvector.
- \rightarrow SELECT ts_rank(

setweight(to_tsvector('english', 'The quick brown fox jumps over the lazy dog'), 'A') setweight(to_tsvector('english', 'An English language pangram. A sentence that contains all of the letters of the alphabet.'), 'B'), plainto_tsquery('english', 'jumping dog')

- → 0.9524299
- different document lengths, should be performed.

The ts rank and ts rank cd functions, return a score for each returned row for a certain match

> You can also change the weights of the tsvector classes (A to D) and set how normalization, due to

Pre-calculate FTS

- is to be performed containing the tsvector values of each row.
- done easily using a trigger.

```
BEGIN
  IF TG_OP = 'INSERT' THEN
  END IF;
  IF TG_OP = 'UPDATE' THEN
      IF NEW.title <> OLD.title THEN
      END IF;
  END IF;
 RETURN NEW;
END
$$ LANGUAGE 'plpgsql';
```

 \rightarrow For performance reasons, we should consider adding a column to tables where FTS

This column should be updated whenever a row changes or is inserted. This can be

CREATE FUNCTION post_search_update() RETURNS TRIGGER AS \$\$ NEW.search = to_tsvector('english', NEW.title); NEW.search = to_tsvector('english', NEW.title);



Indexing FTS

→ To select all posts containing jumping and dog we can use the following query

- → SELECT titleFROM posts WHERE search @@ plainto_tsquery('english', 'jumping dog') ORDER BY ts_rank(search, plainto_tsquery('english', 'jumping dog')) DESC
- \rightarrow To improve the performance of our full text searches, we can use GIN or GiST indexes:
 - \rightarrow CREATE INDEX search_idx ON posts USING GIN (search);
 - → CREATE INDEX search_idx ON posts USING GIST (search);

→ Note that search is a pre-calculated column containing the tsvector of the columns we want to search.

→ Rule of thumb, use GIN if updates to searchable terms are rare and you want to make searches fast.



A6. MediaLibrary Indexes (Performance)

1. Database workload

Understanding the nature of the workload for the application and the performance goals, is essential to develop a good database design. The workload includes an estimate of the number of tuples for each relation and also the estimated growth.

Relation	Relation name	Order of magnitude	Estimated growth
R01	user	10 k (tens of thousands)	10 (tens) / day
R02	author	1 k (thousands)	1 (units) / day
R03	collection	100 (hundreds)	1 / day
R04	work	1 k	1 / day
R05	author_work	1 k	1 / day
R06	nonbook	100	1 / day
R07	publisher	100	1 / day
R08	book	1 k	1 / day
R09	location	100	1 / day
R10	item	10 k	10 / day
R11	loan	100 k	100 (hundreds) / day
R12	review	10 k	10 / day
R13	wish_list	10 k	10 / day

2. Proposed Indexes

Indexes are used to enhance database performance by allowing the database server to find and retrieve specific rows much faster. An index defined on a column that is part of a join condition can also significantly speed up queries with joins. Moreover, indexes can also benefit UPDATE and DELETE commands with search conditions.

After an index is created, the system has to keep it synchronised with the table, which adds overhead to data manipulation operations. As indexes add overhead to the database system as a whole, they are used sensibly.

2.1. Performance indices

Performance indexes are applied to improve the performance of select queries. At most, three performance indexes can be proposed, identifying the ones that have the biggest impact on the performance of the application.

Indexes should be proposed considering queries that are frequently used and involve large relations. Additionally, this section includes an analysis of the execution plan for two central, non-trivial, and frequently used SQL queries significantly impacted by the proposed performance indexes.

Index	IDX01	
Index relation	work	
Index attribute	id_users	
Index type	B-tree	
Cardinality	Medium	
Clustering	Yes	
Justification	Table 'work' is very large. Several queries need to frequently filter access to the works by its owner (user). Filtering is done by exact match, thus an hash type index would be best suited. However, since we also want to apply clustering based on this index, and clustering is not possible on hash type indexes, we opted for a b-tree index. Update frequency is low and cardinality is medium so it's a good candidate for clustering.	
SQL Code		
<pre>CREATE INDEX user_work ON work USING btree (id_users); CLUSTER work USING user_work;</pre>		



A6. MediaLibrary Indexes (Full Text Search)

2.2. Full-text Search indexes

Full-text search indexes are applied to provide keyword based search over records of the database. Results using FTS are ranked by relevance and can use signals from multiple tables and with different weights. The first step in the process of defining FTS indices is to define what is a 'document' for the search features to support.

Index	IDX11
Index relation	work
Index attributes	title, obs
Index type	GIN
Clustering	No
Justification	To provide full-text search features to look for works based on matching titles or observations. The index type is GIN because the indexed fields are not expected to change often.
SQL Code	

-- Add column to work to store computed ts vectors.

SQL Code

```
-- Add column to work to store computed ts_vectors.
ALTER TABLE work
ADD COLUMN tsvectors TSVECTOR;
-- Create a function to automatically update ts_vectors.
CREATE FUNCTION work_search_update() RETURNS TRIGGER AS $$
BEGIN
IF TG_OP = 'INSERT' THEN
       NEW.tsvectors = (
        setweight(to_tsvector('english', NEW.title), 'A') ||
        setweight(to_tsvector('english', NEW.obs), 'B')
        );
 END IF;
 IF TG_OP = 'UPDATE' THEN
        IF (NEW.title <> OLD.title OR NEW.obs <> OLD.obs) THEN
           NEW.tsvectors = (
            setweight(to_tsvector('english', NEW.title), 'A') ||
            setweight(to_tsvector('english', NEW.obs), 'B')
           );
         END IF;
 END IF;
RETURN NEW;
END $$
LANGUAGE plpgsql;
-- Create a trigger before insert or update on work.
CREATE TRIGGER work_search_update
BEFORE INSERT OR UPDATE ON work
 FOR EACH ROW
EXECUTE PROCEDURE work_search_update();
-- Finally, create a GIN index for ts_vectors.
CREATE INDEX search_idx ON work USING GIN (tsvectors);
```



A6. Indexes Checklist

A6. Indexes, Integrity and Populated Database

Artofact	1.1	The artefact reference
Artelact	1.2	The goal of the artef
	2.1	The workload section
Workload	2.2	The relations' magni
	2.3	For each relation, ma
	3.1	Performance indexe
	3.2	For each index, a rel
	3.3	For each index, the t
	3.4	For each index, the
	3.5	For each index, clust
Indovoe	3.6	The impact of the inc
Indexes	3.7	Full-text search (FTS
	3.8	For FTS indexes, fie
	3.9	For each index, a just
	3.10	For each index, the
	3.11	Indexes are not prop
	3.12	Indexed are not prop

ce and name are clear fact is briefly presented (1, 2 sentences) n is included itude and growth estimation section is included agnitude and growth is estimated s are proposed lation and attribute(s) is defined type is defined cardinality is defined tering is defined dices is analysed for two illustrative queries S) indexes over multiple fields are proposed eld weighting is used stification is provided SQL code is included osed for PK posed for UK



A6. Triggers and User Defined Functions

- described by presenting the event, the condition, and the activation code.
- complex computations, are identified and described to be trusted by the database server.
- type or a composite type. Functions can also be defined to return sets of base or composite values.
- \rightarrow Common examples:
 - \rightarrow User cannot post in groups when he is not a member;
 - \rightarrow When a vote is cast, the rating (karma) of the author is updated;
 - \rightarrow When an event happens, relevant notifications are sent.

To enforce integrity rules that cannot be achieved in a simpler way, the necessary triggers are identified and

-> User-defined functions, and trigger procedures, that add control structures to the SQL language, or perform

Every kind of function (SQL functions, Stored procedures, Trigger procedures) can take base types, composite types, or combinations of these as arguments (parameters). In addition, every kind of function can return a base



User-Defined Functions

 \rightarrow A user-defined function provides a mechanism for extending the functionality of the database server by adding a function.

- → Advantages of using stored procedures:
 - Reduce the number of round trips between application and database server
 - uses its full-power
 - → Be able to be reused in many applications
- → Disadvantages of stored procedures:

 - → Make it difficult to manage versions and hard to debug

→ Increase the application performance, because user-defined functions are pre-compiled and stored in the database server

Slow software development because it requires specialized skills that many developers do not possess (PL/SQL)

→ Less portable code to other database management systems (MySQL, SQL Server, PostgreSQL, Oracle, DB2)



UDF Example

```
CREATE [OR REPLACE] FUNCTION function_name (arguments)
RETURNS return_datatype AS $name$
DECLARE
declaration;
[...]
BEGIN
< function_body >
[...]
RETURN { variable_name | VALUE }
END;
$name$ LANGUAGE plpgsql;
```

```
CREATE OR REPLACE FUNCTION totalRecords ()
RETURNS INTEGER AS $total$
DECLARE
  total INTEGER;
BEGIN
  SELECT COUNT(*) INTO total FROM company;
  RETURN total;
END;
$total$ LANGUAGE plpgsql;
```

```
SELECT totalRecords();
```



Iriggers

- → Triggers are <u>event-condition-action</u> rules:
 - \rightarrow Event, a change to the database that activates the trigger
 - \rightarrow Condition, a query or test that is run when the trigger is activated
 - Action, a procedure that is executed when the trigger is activated and its condition is true.
- → The action can be executed *before, after or instead of the trigger* event
- The action may refer the *new values and old values of records inserted*, updated or deleted in the trigger event
- → The programmer specifies that the action is performed:
 - \rightarrow once for each modified record (FOR EACH ROW)
 - \rightarrow once for all records that are changed on a database operation





Triggers Example

CREATE TRIGGER loan_item BEFORE INSERT OR UPDATE ON loan FOR EACH ROW EXECUTE PROCEDURE loan_item();



A6. MediaLibrary Triggers

3. Triggers

Triggers and user defined functions are used to automate tasks depending on changes to the database. Business rules are usually enforced using a combination of triggers and user defined functions.

```
TRIGGER01
Trigger
Description
SQL Code
CREATE FUNCTION loan_item() RETURNS TRIGGER AS
$BODY$
BEGIN
id_users AND end_t > NEW.start_t) THEN
in every moment.';
       END IF;
       RETURN NEW;
END
$BODY$
LANGUAGE plpgsql;
CREATE TRIGGER loan_item
        BEFORE INSERT OR UPDATE ON loan
        FOR EACH ROW
        EXECUTE PROCEDURE loan_item();
```

An item can only be loaned to one user at a time. IF EXISTS (SELECT * FROM loan WHERE NEW.id_users = RAISE EXCEPTION 'An item can only be loaned to one user



A6. Triggers Checklist

Triggers and User Defined Functions	4.1	Triggers and functions	
	4.2	Restrictions not yet co	
	Functions	4.3	For each trigger, a just
		4.4	For each trigger, the S

are proposed
vered in the schema are defined for high priority US
tification is included
QL code is included



A6. Transactions

- \rightarrow For each necessary transaction, include:
 - → Justification
 - → Isolation level
 - \rightarrow SQL code to create it
- \rightarrow Common examples:
 - \rightarrow insert data (e.g. generalizations, latest generated PK)
 - \rightarrow delete data (e.g. user deletes account)
 - \rightarrow checkout purchase (move products from cart to purchase)
 - → many more

-> Transactions bundle multiple steps into a single, all-or-noting operation, ensuring data integrity with concurrent accesses.



Transaction Isolation

- \rightarrow RDBMS offer different isolation levels, which are achieved mostly by locking access to tables.
- \rightarrow Stricter or looser isolation levels will allow less or more concurrent accesses.
- → We should aim to the less restrictive isolation level that still guarantees that data is consistent \rightarrow Advice: declare transactions as READ ONLY when possible.
- \rightarrow Isolation levels are defined in terms of the problems (phenomena) that can occur when concurrent transactions execute:
 - \rightarrow Dirty read a transaction reads data written by a concurrent uncommitted transaction;
 - \rightarrow Non-repeatable read a transaction re-reads data and finds that data has been modified by another transaction;
 - \rightarrow Phantom read a transaction re-executes a query and finds that the results have changed by another transaction;
 - -> Serialization anomaly the result of committing a group of transactions is inconsistent is inconsistent with all possible orderings of running those transactions on at a time.



Transaction Isolation in PostgreSQL

PostgreSQL <u>https://www.postgresql.org/docs/current/transaction-iso.html</u>

- -> serializable transactions see only data committed before the transaction began and never sees uncommitted data or changes
- → repeatable read transactions can only read committed records and between two reads the transactions cannot modify the record
- → read committed transactions can only read committed records (between two reads the record may have been modified)
- \rightarrow read uncommitted records still uncommitted can be read

Table 13.1. Transaction Isolation Levels

Isolation Level	Dirty Read	Nonrepeatable Read	Phantom Read	Serialization Anomaly
Read uncommitted	Allowed, but not in PG	Possible	Possible	Possible
Read committed	Not possible	Possible	Possible	Possible
Repeatable read	Not possible	Not possible	Allowed, but not in PG	Possible
Serializable	Not possible	Not possible	Not possible	Not possible



A6. MediaLibrary Transactions

4. Transactions

Transactions are used to assure the integrity of the data when multiple operations are necessary.

Transaction	TRAN01	
Description	Get current	
Justification	In the midd the loan tab retrieved in a Phantom Selects.	
Isolation level	SERIALIZA	
SQL Code		
BEGIN TRANSACTION	N;	
SET TRANSACTION	ISOLATION I	
Get number of SELECT COUNT(*) FROM loan WHERE now() < end	current lo d_t;	
Get ending loa SELECT loan.end_t users.name FROM loan INNER JOIN item (INNER JOIN work (INNER JOIN users WHERE now () < lo	ans (limit t, loan.sta ON item.id ON work.id ON users.a	
ORDER BY loan.end LIMIT 10;	d_t ASC	
END TRANSACTION;		

Ile of the transaction, the insertion of new rows in ble can occur, which implies that the information both selects is different, consequently resulting in Read. It's READ ONLY because it only uses

LEVEL SERIALIZABLE READ ONLY;

oans

```
10)
art_t, item.*, work.*, users.id,
```

```
= loan.id_item
= item.id_work
id = loan.id_users
```



A6. Transactions Checklist

Database transactions	5.1	Database transactions
	5.2	Each transaction has a
	5.3	Each transaction has a
	5.4	Transactions' SQL syn
	5.5	No unnecessary transa
	5.6	All transactions for hig

section is included			
n isolation type defined and justified			
justification			
ax is correct			
ctions are included			
priority users stories are included			



A6. Database Population

- The EBD Component includes two SQL scripts
 - functions;
 - plausible values for each field type.
- \rightarrow These scripts must run, *as-is*, in the production PostgreSQL environment.

→ Database creation script, including SQL creation statements for all tables, key constraints, performance indexes, full text search indexes, triggers, user defined

Database population script, including SQL insert statements to populate a database with test data with an amount of tuples suitable for testing and with



A6. MediaLibrary Database Population

A.1 Database schema

```
    Drop old schema

DROP TABLE IF EXISTS wish_list CASCADE;
DROP TABLE IF EXISTS review CASCADE;
DROP TABLE IF EXISTS loan CASCADE;
DROP TABLE IF EXISTS item CASCADE;
DROP TABLE IF EXISTS nonbook CASCADE;
DROP TABLE IF EXISTS book CASCADE;
DROP TABLE IF EXISTS author_work CASCADE;
DROP TABLE IF EXISTS work CASCADE;
DROP TABLE IF EXISTS collection CASCADE;
DROP TABLE IF EXISTS author CASCADE;
DROP TABLE IF EXISTS location CASCADE;
DROP TABLE IF EXISTS publisher CASCADE;
DROP TABLE IF EXISTS users CASCADE;
DROP TYPE IF EXISTS media;
 - Types
 ______
CREATE TYPE media AS ENUM ('CD', 'DVD', 'VHS', 'Slides', 'Photos',
'MP3');
 ______
 - Tables
 Note that a nlural 'users' name was adopted because user is a
```

A.2 Database population



A6. Database Creation and Population Checklist

SQL	6.1	The SQL schema scri
	6.2	The SQL script resets
	6.3	The SQL schema scri
	6.4	The SQL population s
	6.5	The SQL population s
	6.6	The SQL population s
	6.7	The SQL schema scri
	6.8	The production databa

pt is included

the database state (includes DROPs + CREATEs)

pt executes without errors

cript is included

cript is included in the group's repository

cript executes without errors

pt is included in the group's repository

ase (at db.fe.up.pt) has been updated with the SQL scripts



PostgreSQL

Docker

- Docker is a key technology in LBAW (PostgreSQL, pgAdmin, Laravel)
- \rightarrow It is mandatory for deploying your prototypes and final products
- Docker is a lightweight virtualization environment, widely used to package applications and its dependencies in isolated containers.
- \rightarrow With Docker you can manage your product infrastructure as applications.
- Available for Windows, Mac and Linux <u>https://docs.docker.com/get-docker/</u>
- → Important: don't postpone using Docker.



PostgreSQL Docker Container

- Official PostgresSQL are available at: <u>https://hub.docker.com/ /postgres</u>
- → Start a local PostgreSQL server with:
 - → docker run --name some-postgres -e POSTGRES_PASSWORD=mysecretpassword -d postgres
- \rightarrow Run a local pgAdmin installation (available at localhost:80) with:
 - \rightarrow docker run -p 80:80 \
 - -e 'PGADMIN_DEFAULT_EMAIL=user@domain.com' \
 - -e 'PGADMIN_DEFAULT_PASSWORD=SuperSecret' \
 - -d dpage/pgadmin4



Docker Compose

- → Docker Compose is used to setup multi-container Docker applications.
- \rightarrow A YAML file is used to configure the containers to start and run.
- The LBAW 'template-postgresql' repository, setups up two containers <u>https://git.fe.up.pt/lbaw/template-postgresql</u>

```
version: '3'
services:
 postgres:
    image: postgres:11.13
   restart: always
    environment:
     POSTGRES_USER: postgres
     POSTGRES PASSWORD: pg!password
   ports:
      - "5432:5432"
 pgadmin:
   image: dpage/pgadmin4:6
    environment:
    ports:
      - "4321:80"
   depends_on:
      - postgres
```

PGADMIN_DEFAULT_EMAIL: postgres@lbaw.com PGADMIN_DEFAULT_PASSWORD: pg!password



About the PostgreSQL Production Environment (important!)

- \rightarrow A PostgreSQL database contains one or more schemas, which in turn contains one or more tables.
- \rightarrow All databases contain a public schema, which is used as default.
- → In PostgreSQL's command line interface, you can view the current active schema with: show search_path;
- \rightarrow To change the schema for the current session use: SET search_path TO <schema>;

\rightarrow In the PostgreSQL setup at FEUP (db.fe.up.pt), the public schema is shared between all accounts,

 \rightarrow Tables created in the public schema are visible to all users (although not accessible). If you look at the tables in the publish schema, you will find a long list of tables.

→ It is important to not use the public schema and instead create a schema with the name of your group (lbaw21gg).

- \rightarrow To create this schema, use the following command: CREATE SCHEMA < 1baw21gg>;
- → To always use this schema as the default in your project, add the following line to the beginning of your SQL scripts.

→ SET search_path TO <lbaw21gg>;









References

Bibliography and Further Reading

- PostgreSQL Manual, Chapter 11. Indexes, <u>https://www.postgresql.org/docs/14/indexes.html</u>
- PostgreSQL Manual, Chapter 12. Full Text Search, <u>https://www.postgresql.org/docs/14/textsearch.html</u>
- → Scott Ambler, The Object Primer, Cambridge University Press, 3rd Edition, 2004.
- → Alberto Rodrigues da Silva, Carlos Videira, UML Metodologias e Ferramentas CASE, 2ª Edição, Centro Atlântico Editora, Maio 2005.
- Raghu Ramakrishnan, Johannes Gehrke. Database Management Systems. McGRAW-Hill International Editions, 3rd Edition, 2003.



Lab Class #4

- \rightarrow Discuss the conceptual data model (A4)
- \rightarrow Develop and discuss the relational schema (A5)
 - → Map the classes and relationships of the conceptual schema into relation schemas
 - \rightarrow For each relation, identify the functional dependencies (FD) that apply
 - → Check if each relation is in BCNF

 - 3NF, preserving the functional dependencies
 - → Develop and test a first version of the database creation script in SQL

- Test the local development environment for PostgreSQL.
- \rightarrow Test the connection to the production PostgreSQL server at <u>db.fe.up.pt</u>

→ If the relation is not in BCNF and there are no other impediments, look for several possible decompositions (lossless)

> If there is no satisfactory decomposition to BCNF and if the relation is no longer in 3NF, consider the decomposition lossless for

